NAME:	
COLLABORATOR(S):_	

THIS IS AN OPTIONAL HOMEWORK TO REPLACE A PREVIOUS HOMEWORK GRADE

1/2/1/0	1. Explain why the following code sr	nippet is not atomic?	
5/3/1/0	balance = ba	alance + 1	
3/5/3/0	2. In the following code snippet what program? Is the expected output consthe program? Explain?		
int s	hared;		
i f } r } int m p p	<pre>* fun(void * args){ nt i; or(i=0;i<100;i++){ shared++; return NULL; ain(){ thread_t t1,t2; thread_create(&t1, NULL, fun, NULL); thread_create(&t2, NULL, fun, NULL); thread_join(t1, NULL); thread_join(t2, NULL); orintf("shared: %d\n", shared);_</pre>		
}			
5/3/1/0	3. In the above code snippet circle explain describe a critical section.		

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4. Consider the naïve locking solution used for the thread startup routine from the previous program: Does this provide proper locking? Why or why not, explain.

```
int shared;
int lock;
void * fun(void * args){
   int i;

   for(i=0;i<100;i++){
      while(lock > 0);//spin
      lock = 1; //set lock
      shared++; //increment
      lock = 0; //unlock
   }

   return NULL;
}
```

5/3/1/0 5. Explain why using a **mutex** avoids issues of a lack of atomicity in lock acquisition?

7/5/3/0 6. Which type of locking strategy, coarse or fine, does the following code block use? Is there a possibility of a more efficient locking strategy? Explain.

```
pthread_mutext_t lock;
int avail = MAX_FUNDS;
int local 1 = 0;
int local 2 = 0;
void * fun(void * args){
   int v,i;
    for(i=0; i < 100; i++){
        v = random() % 100;
        pthread_mutext_lock(&lock);
        if(avail - v > 0){
            avail -= v;
        if(random() % 2){
            local_1 += v;
        }else{
            local_2 += v;
        }
        pthread_mutext_unlock(&lock);
    return NULL;
```

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10/8/6/3/0

7. Based on the code example from Question 6, fill in locking code to provide a more efficient locking strategy.

```
int avail = MAX_FUNDS;
int local_1 = 0;
int local 2 = 0;
void * fun(void * args){
    int v,i;
    for(i=0; i < 100; i++){
        v = random() % 100;
        if(avail - v > 0){
            avail -= v;
        }
        if(random() % 2){
            local_1 += v;
        }else{
            local_2 += v;
       }
    return NULL:
```

5/3/1/0 8. What is deadlock and provide a small (pseudo-)code example of how deadlock can arrise from coarse grain locking.

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5/3/1/0 9. Provide an example of deadlock avoidance when there is a natural ordering of lockable objects.

10/8/6/3/0	10. Provide a detailed description of the problem setup for the dining philosophers problem:
	diffing philosophers problem.
25/23/20/	15/10/5/0
	11. In pseudo code, provie a solution to the dining philosophers problem that avoids deadlock:
10/8/6/3/0	12. Explain your solution and argue that it wll always avoid deadlocks regardless of the number of philosophers.

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